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SOME GENERALIZATIONS OF p-LOXODROMIC FUNCTIONS

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Denote $\mathbb{C}^* = \mathbb{C} \setminus \{0\}$ and let $q, p \in \mathbb{C}^*$, |q| < 1.

Definition. [1] A meromorphic in \mathbb{C}^* function f is said to be p-loxodromic of multiplicator q if for every $z \in \mathbb{C}^*$

$$f(qz)=pf(z).$$

For $z \in \mathbb{C}^*$ consider the equation of the form

$$f(qz) = p(z)f(z), (1)$$

where p(z) is some function. If $p(z) \equiv const$, then meromorphic solution of this equation is p-loxodromic function. In particular, if $p(z) \equiv 1$, we have classic loxodromic function. It was studied in the works of O. Rausenberger [2], G. Valiron [3] and Y. Hellegouarch [4]. The class of loxodromic functions is denoted by \mathcal{L}_q .

For certain functions p(z) holomorphic solutions of equation (1) are found. These solutions will be some generalizations of p-loxodromic functions.

First, consider the functional equation of the form

$$f(qz) = \frac{1}{1-z}f(z), z \in \mathbb{C}^*.$$
 (2)

Define the entire function with the zero sequence $\{q^{-n}\}$, where $n \in \mathbb{N} \cup \{0\}$, 0 < |q| < 1,

$$H(z) = \prod_{n=0}^{\infty} (1 - q^n z).$$

Theorem 1. Every holomorphic in \mathbb{C}^* solution of (2) has the form f(z) = CH(z), where C is a constant.

Also, we consider the functional equation of the form

$$f(qz) = \frac{1}{z}f(z), z \in \mathbb{C}^*.$$
 (3)

Definition. |1| The function

$$P(z) = (1-z) \prod_{n=1}^{\infty} (1-q^n z) \left(1 - \frac{q^n}{z}\right)$$

is called the Schottky-Klein prime function.

Theorem 2. Every holomorphic in \mathbb{C}^* solution of (3) has the form f(z) = CP(-z), where C is a constant.

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Category of ambiguous representations

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It is well known since the pioneering work [1] that a computer program is described by the corresponding predicate transformer, i.e. the mapping that transforms a valid knowledge about program input into a valid knowledge about the output. It is also commonly agreed [3] that the proper domains containing such predicates are directed complete continuous posets, and the predicate transformers are Scott continuous. Hence a viable choice for description of deterministic programs is the category $\mathcal{S}\text{em}_0$ of all continuous semigroups [2] with bottom elements and their Scott continuous bottom-preserving (not necessarily meet-preserving) mappings.

Unfortunately this category is inappropriate for *nondeterministic* programs, or, equivalently, for game situations, where a player can play different moves in the same position. We propose to use a relation (= a multivalued mapping) here: